Probabilistic & Unsupervised Learning

Convex Algorithms in Approximate Inference

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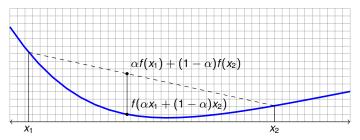
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Convexity

A convex function $f : X \to \mathbb{R}$ is one where

 $f(\alpha x_1 + (1 - \alpha)x_2) \leq \alpha f(x_1) + (1 - \alpha)f(x_2)$

for any $x_1, x_2 \in X$ and $0 \le \alpha \le 1$.



Convex functions have a global infimum (unless not bounded below) and there are efficient algorithms to find a minimum subject to convex constraints.

Examples: linear programs (LP), quadratic programs (QP), second-order cone programs (SOCP), semi-definite programs (SDP), geometric programs.

Convexity and Approximate Inference

The theory of convex functions and convex spaces has long been central to optimisation. It has recently also found application in the theory of free energy and approximation:

- Linear programming relaxation as an approximate method to find the MAP assignment in Markov random fields.
- Attractive Markov random fields: binary case exact and related to a maximum flow-minimum cut problem in graph theory (a linear program). Approximate otherwise.
- Unified view of approximate inference as optimization on the marginal polytope.
- Tree-structured convex upper bounds on the log partition function (convexified belief propagation).
- Learning graphical models using maximum margin principles and convex approximate inference.

Consider a discrete Markov random field (MRF) with pairwise interactions:

$$p(\mathbf{X}) = \frac{1}{Z} \prod_{(ij)} f_{ij}(X_i, X_j) \prod_i f_i(X_i) = \frac{1}{Z} \exp\left(\sum_{(ij)} E_{ij}(X_i, X_j) + \sum_i E_i(X_i)\right)$$

The problem is to find the most likely configuration \mathbf{X}^{MAP} :

$$\mathbf{X}^{\mathsf{MAP}} = \operatorname*{argmax}_{\mathbf{X}} \sum_{(ij)} E_{ij}(X_i, X_j) + \sum_i E_i(X_i)$$

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Reformulate in terms of indicator variables:

$$b_i(k) = \delta(X_i = k)$$

$$b_{ij}(k, l) = \delta(X_i = k)\delta(X_j = l)$$

where $\delta(\cdot) = 1$ if argument is true, 0 otherwise. Each $b_i(k)$ is an indicator for whether variable X_i takes on value k. The indicator variables need to satisfy certain constraints:

$$\begin{split} b_i(k), b_{ij}(k,l) &\in \{0,1\} & \text{Indicator variables are binary variables.} \\ &\sum_k b_i(k) = 1 & X_i \text{ takes on exactly one value.} \\ &\sum_l b_{ij}(k,l) = b_i(k) & \text{Pairwise indicators are consistent with single-site indicators.} \end{split}$$

MAP assignment problem is equivalent to:

$$\underset{\{b_i, b_{ij}\}}{\operatorname{argmax}} \sum_{(ij)} \sum_{k, l} b_{ij}(k, l) E_{ij}(k, l) + \sum_i \sum_k b_i(k) E_i(k)$$

with constraints:

$$\forall i, j, k, l: b_i(k), b_{ij}(k, l) \in \{0, 1\}$$
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The linear programming relaxation for MRFs is:

$$\operatorname*{argmax}_{\{b_i,b_{ij}\}}\sum_{(ij)}\sum_{k,l}b_{ij}(k,l)E_{ij}(k,l)+\sum_i\sum_kb_i(k)E_i(k)$$

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$$\forall i, j, k, l:$$
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- LP relaxation is a zero-temperature version of the Bethe free energy formulation of loopy BP, where the Bethe entropy term can be ignored.
- If the MRF is binary and attractive, then (a slightly different reformulation of LP relaxation) will always give the MAP solution.
- Next: we show how to find the MAP solution directly for binary attractive MRFs using network flow.

Binary MRFs:

$$\rho(\mathbf{X}) = \frac{1}{Z} \exp\left(\sum_{(ij)} W_{ij} \delta(X_i = X_j) + \sum_i c_i X_i\right)$$

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$$p(\mathbf{X}) = rac{1}{Z} \exp\left(\sum_{(ij)} W_{ij} \delta(X_i = X_j) + \sum_i c_i X_i\right)$$

The binary MRF is attractive if $W_{ij} \ge 0$ for all i, j.

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- Many practical MRFs are attractive, e.g. image segmentation, webpage classification.
- MAP X can be found efficiently by converting problem into a maximum flow-minimum cut program.

The MAP problem:

$$\operatorname*{argmax}_{\mathbf{x}}\sum_{(ij)} \textit{W}_{ij}\delta(\textit{x}_i=\textit{x}_j) + \sum_i \textit{c}_i\textit{x}_i$$

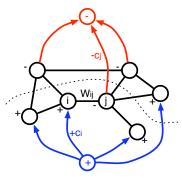
Construct a network as follows:

- 1. Edges (*ij*) are undirected with weight $\lambda_{ij} = W_{ij}$;
- 2. Add a source *s* and a sink *t* node;
- 3. $c_i > 0$: Connect the source node to variable *i* with weight $\lambda_{si} = c_i$;
- 4. $c_j < 0$: Connect variable *j* to the sink node with weight $\lambda_{jt} = -c_j$.

A cut is a partition of the nodes into S and T with $s \in S$ and $t \in T$. The weight of the cut is

$$\Lambda(S,T) = \sum_{i \in S, j \in T} \lambda_{ij}$$

The minimum cut problem is to find the cut with minimum weight.



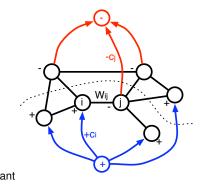
Identify an assignment $\mathbf{X} = \mathbf{x}$ with a cut:

 $S = \{s\} \cup \{i : x_i = 1\}$ $T = \{t\} \cup \{j : x_j = 0\}$

The weight of the cut is:

Λ

$$egin{aligned} \mathcal{M}(\mathcal{S},\mathcal{T}) &= \sum_{(ij)} \mathcal{W}_{ij}\delta(x_i
eq x_j) \ &+ \sum_i (1-x_i)\max(0,c_i) \ &+ \sum_j x_j\max(0,-c_j) \ &= -\sum_{(ij)} \mathcal{W}_{ij}\delta(x_i=x_j) - \sum_i x_ic_i + ext{constants} \end{aligned}$$



So finding the minimum cut corresponds to finding the MAP assignment.

How do we find the minimum cut? The minimum cut problem is dual to the **maximum flow problem**, i.e. find the maximum flow allowable from the source to the sink through the network. This can be solved extremely efficiently (see wikipedia entry).

The framework can be generalized to general attractive MRFs, but will not be exact anymore.

Convexity in exponential family inference and learning

Exponential families: the log partition function

Consider an exponential family distribution with sufficient statistic s(X) and natural parameter θ (and no base factor in *X* alone). We can write its probability or density function as

$$p(X|\theta) = \exp\left(\theta^{\mathsf{T}}s(X) - \Phi(\theta)\right)$$

where $\Phi(\theta)$ is the log partition function

$$\Phi(\boldsymbol{\theta}) = \log \sum_{x} \exp\left(\boldsymbol{\theta}^{\mathsf{T}} \boldsymbol{s}(x)\right)$$

 $\Phi(\theta)$ plays an important role in the theory of the exponential family. For example, it maps natural parameters to the moments of the sufficient statistics:

$$\frac{\partial}{\partial \theta} \Phi(\theta) = e^{-\Phi(\theta)} \sum_{x} s(x) e^{\theta^{\mathsf{T}} s(x)} = \mathbb{E}_{\theta} [s(X)] = \mu(\theta) = \mu$$
$$\frac{\partial^{2}}{\partial \theta^{2}} \Phi(\theta) = e^{-\Phi(\theta)} \sum_{x} s(x)^{2} e^{\theta^{\mathsf{T}} s(x)} - e^{-2\Phi(\theta)} \left[\sum_{x} s(x) e^{\theta^{\mathsf{T}} s(x)} \right]^{2} = \mathbb{V}_{\theta} [s(X)]$$

The second derivative is thus positive semi-definite, and so $\Phi(\theta)$ is convex in θ .

Exponential families: mean parameters and negative entropy

An exponential family distribution can also (almost always) be parameterised by the means of the sufficient statistics.

 $\mu(\theta) = \mathbb{E}_{\theta}\left[s(X)\right]$

Consider the negative entropy of the distribution as a function of the mean parameter:

$$\Psi(\boldsymbol{\mu}) = \mathbb{E}_{\boldsymbol{\theta}} \left[\log p(\boldsymbol{X} | \boldsymbol{\theta}(\boldsymbol{\mu})) \right] = \boldsymbol{\theta}^{\mathsf{T}} \boldsymbol{\mu} - \Phi(\boldsymbol{\theta})$$

so

$$\boldsymbol{\theta}^{\mathsf{T}} \boldsymbol{\mu} = \Phi(\boldsymbol{\theta}) + \Psi(\boldsymbol{\mu})$$

The negative entropy is dual to the log-partition function. For example,

$$egin{aligned} &rac{d}{d\mu}\Psi(\mu) = rac{\partial}{\partial\mu}ig(heta^{ extsf{ iny T}}\mu - \Phi(heta)ig) + rac{d heta}{d\mu}rac{\partial}{\partial heta}ig(heta^{ extsf{ iny T}}\mu - \Phi(heta)ig) \ &= heta + rac{d heta}{d\mu}(\mu-\mu) = heta \end{aligned}$$

Exponential families: duality

In fact, the log partition function and negative entropy are conjugate dual functions.

Consider the KL divergence between distributions with natural parameters θ and θ' :

$$\begin{split} \mathsf{KL}\big[\theta\big\|\theta'\big] &= \mathsf{KL}\big[p(X|\theta)\big\|p(X|\theta')\big] = \mathbb{E}_{\theta}\left[-\log p(X|\theta') + \log p(X|\theta)\right] \\ &= -\theta'^{\mathsf{T}}\mu + \Phi(\theta') + \Psi(\mu) \geq 0 \\ &\Rightarrow \Psi(\mu) \geq \theta'^{\mathsf{T}}\mu - \Phi(\theta') \end{split}$$

where μ are the mean parameters corresponding to θ .

Now, the minimum KL divergence of zero is reached iff $oldsymbol{ heta}=oldsymbol{ heta}'$, so

$$\Psi(\mu) = \sup_{\theta'} \left[\theta'^{\mathsf{T}} \mu - \Phi(\theta') \right] \qquad \text{and, if finite} \quad \theta(\mu) = \operatorname*{argmax}_{\theta'} \left[\theta'^{\mathsf{T}} \mu - \Phi(\theta') \right]$$

The left-hand equation is the definition of the conjugate dual of a convex function.

Continuous functions are reciprocally dual, so we also have:

$$\Phi(\theta) = \sup_{\mu'} \left[\theta^{\mathsf{T}} \mu' - \Psi(\mu') \right] \qquad \text{and, if finite} \quad \mu(\theta) = \underset{\mu'}{\operatorname{argmax}} \left[\theta^{\mathsf{T}} \mu' - \Psi(\mu') \right]$$

Thus, duality gives us another relation between θ and μ .

Duality, inference and the free energy

Consider a joint exponential family distribution on observed x and latent y.

$$p(\mathbf{x}, \mathbf{y}) = \exp\left[\mathbf{\theta}^{\mathsf{T}} s(\mathbf{x}, \mathbf{y}) - \Phi_{XY}(\mathbf{\theta})
ight]$$

The posterior on **y** is also in the exponential family, with the clamped sufficient statistic $s_Y(\mathbf{y}; \mathbf{x}) = s_{XY}(\mathbf{x}^{\text{obs}}, \mathbf{y})$; the same (now possibly redundant) natural parameter θ ; and partition function $\Phi_Y(\theta) = \log \sum_{\mathbf{y}} \exp \theta^T s_Y(\mathbf{y})$.

The likelihood is

$$\mathcal{L}(\theta) = p(\mathbf{x}|\theta) = \sum_{\mathbf{y}} e^{\theta^{\mathsf{T}} s(\mathbf{x}, \mathbf{y}) - \Phi_{XY}(\theta)} = \sum_{\mathbf{y}} e^{\theta^{\mathsf{T}} s_{Y}(\mathbf{y}; \mathbf{x})} e^{-\Phi_{XY}(\theta)} = \exp[\Phi_{Y}(\theta) - \Phi_{XY}(\theta)]$$

So we can write the log-likelihood as

$$\ell(\boldsymbol{\theta}) = \sup_{\boldsymbol{\mu}_{\boldsymbol{Y}}} [\underbrace{\boldsymbol{\theta}^{\mathsf{T}} \boldsymbol{\mu}_{\boldsymbol{Y}} - \boldsymbol{\Phi}_{\boldsymbol{X}\boldsymbol{Y}}(\boldsymbol{\theta})}_{\langle \log \boldsymbol{\rho}(\mathbf{x}, \mathbf{y}) \rangle_{q}} - \underbrace{\boldsymbol{\Psi}(\boldsymbol{\mu}_{\boldsymbol{Y}})}_{-\mathbf{H}[q]}] = \sup_{\boldsymbol{\mu}_{\boldsymbol{Y}}} \mathcal{F}(\boldsymbol{\theta}, \boldsymbol{\mu}_{\boldsymbol{Y}})$$

This is the familiar free energy with $q(\mathbf{y})$ represented by its mean parameters μ_{γ} !

We have described inference in terms of the distribution q, approximating as needed, then computing expected suff stats. Can we describe it instead as an optimisation over μ directly?

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Concave maximisation(!), but two complications:

The optimum must be found over feasible means. Interdependance of the sufficient statistics may prevent arbitrary sets of mean sufficient statistics being achieved

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 - Feasible means are convex combinations of all the single-configuration sufficient statistics.

$$\mu = \sum_{\mathbf{x}} \nu(\mathbf{x}) s(\mathbf{x}) \qquad \sum_{\mathbf{x}} \nu(\mathbf{x}) = 1$$

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- Take a Boltzmann machine on two variables, x₁, x₂.
- The sufficient stats are $s(\mathbf{x}) = [x_1, x_2, x_1x_2]$.
- Clearly only the stats $S = \{[0, 0, 0], [0, 1, 0], [1, 0, 0], [1, 1, 1]\}$ are possible.
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- Thus $\mu \in \text{convex hull}(S)$.
- For a discrete distribution, this space of possible means is bounded by exponentially many hyperplanes connecting the discrete configuration stats: called the marginal polytope.
- Even when restricted to the marginal polytope, evaluating $\Psi(\mu)$ can be challenging.

Convexity and undirected trees

• We can parametrise a discrete pairwise MRF as follows:

$$p(\mathbf{X}) = \frac{1}{Z} \prod_{i} f_{i}(X) \prod_{(ij)} f_{ij}(X_{i}, X_{j})$$
$$= \exp\left(\sum_{i} \sum_{k} \theta_{i}(k) \delta(X_{i} = k) + \sum_{(ij)} \sum_{k,l} \theta_{ij}(k, l) \delta(X_{i} = k) \delta(X_{j} = l) - \Phi(\theta)\right)$$

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So discrete MRFs are always exponential family, with natural and mean parameters:

$$\begin{aligned} \boldsymbol{\theta} &= \begin{bmatrix} \boldsymbol{\theta}_i(k), \boldsymbol{\theta}_{ij}(k,l) & \forall i, j, k, l \end{bmatrix} \\ \boldsymbol{\mu} &= \begin{bmatrix} \boldsymbol{p}(X_i = k), \boldsymbol{p}(X_i = k, X_j = l) & \forall i, j, k, l \end{bmatrix} \end{aligned}$$

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If the MRF has tree structure T, the negative entropy can be written in terms of the single-site entropies and mutual informations on edges:

$$\Psi(\boldsymbol{\mu}_{T}) = \mathbb{E}_{\boldsymbol{\theta}_{T}} \left[\log \prod_{i} \boldsymbol{p}(X_{i}) \prod_{(ij) \in T} \frac{\boldsymbol{p}(X_{i}, X_{j})}{\boldsymbol{p}(X_{i}) \boldsymbol{p}(X_{j})} \right]$$
$$= -\sum_{i} H(X_{i}) + \sum_{(ij) \in T} I(X_{i}, X_{j})$$

We can see the Bethe free energy problem as a relaxation of the true free-energy optimisation:

$$\boldsymbol{\mu}_{Y}^{*} = \operatorname*{argmax}_{\boldsymbol{\mu}_{Y} \in \mathcal{M}} [\boldsymbol{ heta}^{\mathsf{T}} \boldsymbol{\mu}_{Y} - \Psi(\boldsymbol{\mu}_{Y})]$$

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 \mathcal{L} is still a convex set (polytope for discrete problems). However Ψ_{Bethe} is not convex.

Convexifying BP

Consider instead an upper bound on $\Phi(\theta)$:

Imagine a set of spanning trees *T* for the MRF, each with its own parameters θ_T , μ_T . By padding entries corresponding to off-tree edges with zero, we can assume that θ_T has the same dimensionality as θ .

Suppose also that we have a distribution β over the spanning trees so that $\mathbb{E}_{\beta} [\boldsymbol{\theta}_{T}] = \boldsymbol{\theta}$.

Then by the convexity of $\Phi(\theta)$,

$$\Phi(oldsymbol{ heta}) = \Phi(\mathbb{E}_eta \left[oldsymbol{ heta}_ oldsymbol{ heta}
ight]) \leq \mathbb{E}_eta \left[\Phi(oldsymbol{ heta}_ oldsymbol{ heta})
ight]$$

If we were to tighten the upper bound we might obtain a good approximation to Φ :

$$\Phi(\boldsymbol{\theta}) \leq \inf_{\beta, \boldsymbol{\theta}_{\mathcal{T}}: \mathbb{E}_{\beta}[\boldsymbol{\theta}_{\mathcal{T}}] = \boldsymbol{\theta}} \mathbb{E}_{\beta} \left[\Phi(\boldsymbol{\theta}_{\mathcal{T}}) \right]$$

Convex Upper Bounds on the Log Partition Function

$$\Phi(\boldsymbol{\theta}) \leq \inf_{\boldsymbol{\theta}_{\mathcal{T}}: \mathbb{E}_{\beta}[\boldsymbol{\theta}_{\mathcal{T}}] = \boldsymbol{\theta}} \mathbb{E}_{\beta} \left[\Phi(\boldsymbol{\theta}_{\mathcal{T}}) \right]$$

Solve this constrained optimisation problem using Lagrange multipliers:

$$\mathcal{L} = \mathbb{E}_{eta} \left[\Phi(oldsymbol{ heta}_{ au})
ight] - oldsymbol{\lambda}^{\mathsf{T}} (\mathbb{E}_{eta} \left[oldsymbol{ heta}_{ au}
ight] - oldsymbol{ heta})$$

Setting the derivatives wrt θ_T to zero, we get:

$$eta(au)oldsymbol{\lambda}_{ au}-eta(au) {\Pi}_{ au}(oldsymbol{\lambda})= 0 \ oldsymbol{\lambda}_{ au}= {\Pi}_{ au}(oldsymbol{\lambda})$$

where $\Pi_{\tau}(\lambda)$ are the Lagrange multipliers corresponding to vertices and edges on the tree τ . Although there can be many θ_{τ} parameters, at optimum they are all constrained: their corresponding mean parameters are all consistent with each other and with λ .

Convex Upper Bounds on the Log Partition Function

$$egin{aligned} \Phi(m{ heta}) &\leq \sup_{m{\lambda}} \inf_{m{ heta}_{ au}} \mathbb{E}_{m{eta}} \left[\Phi(m{ heta}_{ au})
ight] - m{\lambda}^{\mathsf{T}} (\mathbb{E}_{m{eta}} \left[m{ heta}_{ au}
ight] - m{ heta}) \ &= \sup_{m{\lambda}} m{\lambda}^{\mathsf{T}} m{ heta} + \mathbb{E}_{m{eta}} \left[\inf_{m{ heta}_{ au}} \Phi(m{ heta}_{ au}) - m{ heta}_{ au}^{\mathsf{T}} \Pi_{ au}(m{\lambda})
ight] \ &= \sup_{m{\lambda}} m{\lambda}^{\mathsf{T}} m{ heta} + \mathbb{E}_{m{eta}} \left[-\Psi(\Pi_{ au}(m{\lambda}))
ight] \ &= \sup_{m{\lambda}} m{\lambda}^{\mathsf{T}} m{ heta} + \mathbb{E}_{m{eta}} \left[\sum_{i} H_{m{\lambda}}(X_{i}) - \sum_{(ij)\in\mathcal{T}} I_{m{\lambda}}(X_{i}, X_{j})
ight] \ &= \sup_{m{\lambda}} m{\lambda}^{\mathsf{T}} m{ heta} + \sum_{i} H_{m{\lambda}}(X_{i}) - \sum_{(ij)} m{eta}_{ij} I_{m{\lambda}}(X_{i}, X_{j}) \end{aligned}$$

This is a **convexified** Bethe free energy.

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